

# Approximation Algorithms for Continuous Clustering and Facility Location Problems

Deeparnab Chakrabarty, Maryam Negahbani, **Ankita Sarkar**

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  - $k$ -median:  $\sum_{v \in C} d(v, S)$
  - $k$ -means:  $\sum_{v \in C} d(v, S)^2$

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Hochbaum and Shmoys, 1985

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- Discrete best:

- $k$ -median: 2.675

[Byrka, Pensyl, Rybicki, Srinivasan, and Trinh, TALG 2017]

- $k$ -means: 9

[Kanungo, Mount, Netanyahu, Piatko, Silverman, and Wu, ComGeo 2004] ...

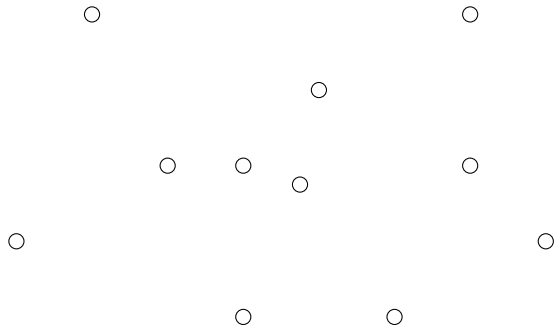


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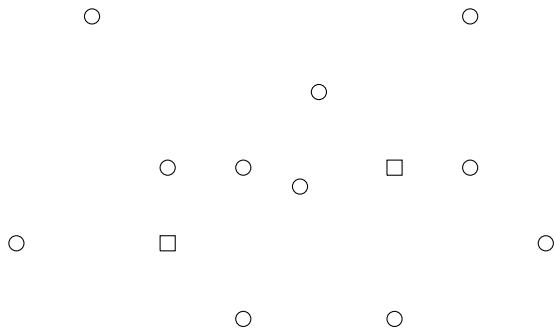
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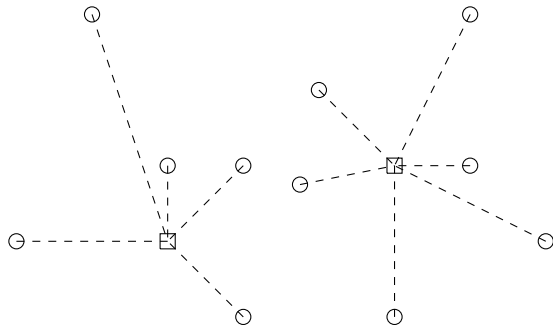
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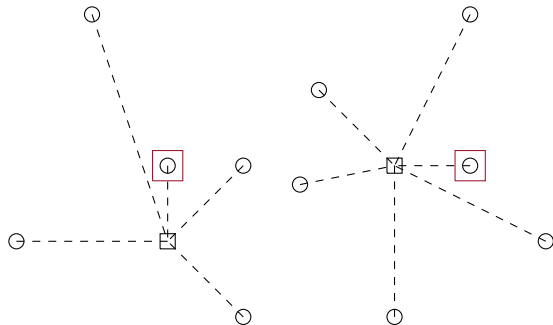
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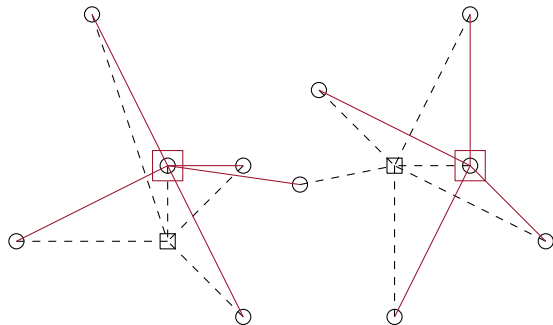
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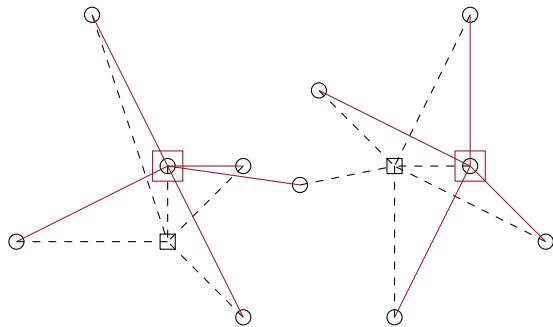
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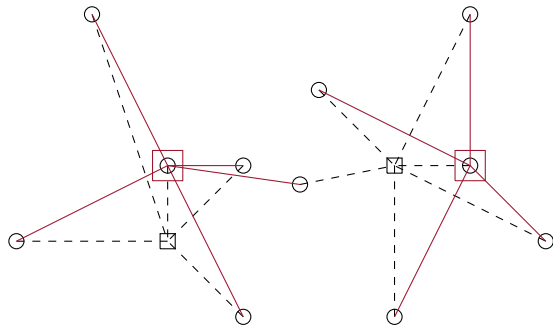
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- General  $(X, d)$ ?

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- Fair- $k$ -median,  $k$ -center with outliers : **Yes.**

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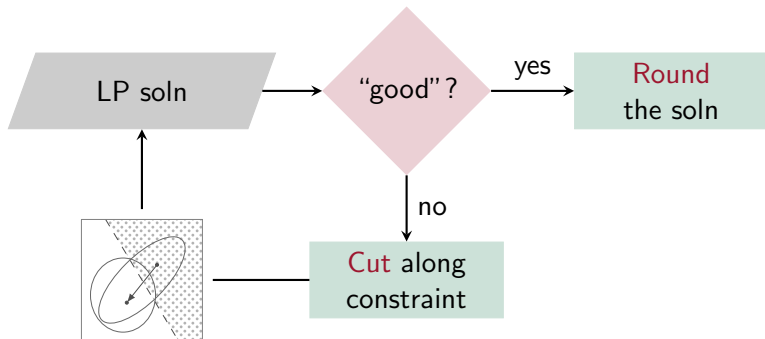
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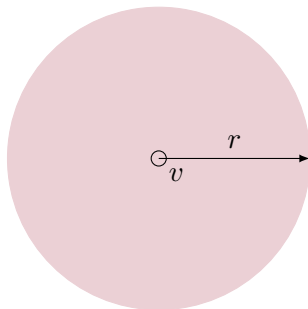
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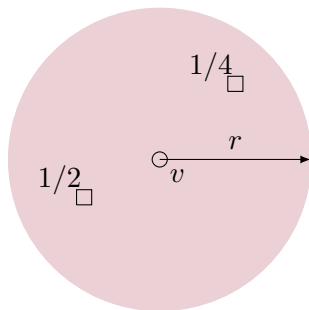


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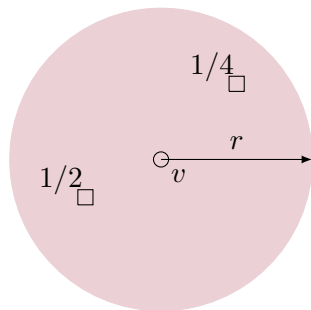


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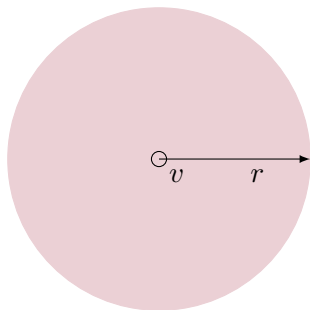


$$y(v, r) = 3/4$$

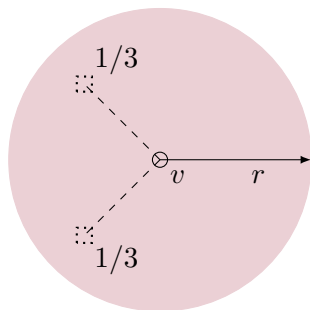
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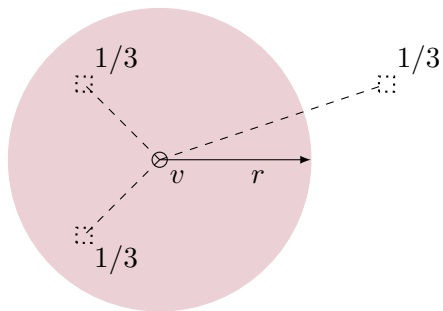


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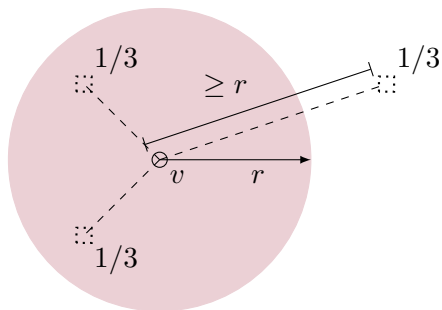
$$\underbrace{y(v, r)}_{2/3}$$

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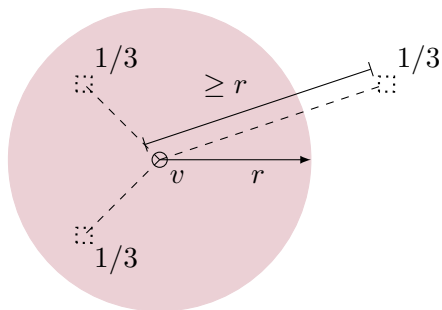
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$$C_v \geq r \underbrace{(1 - y(v, r))}_{1/3}$$

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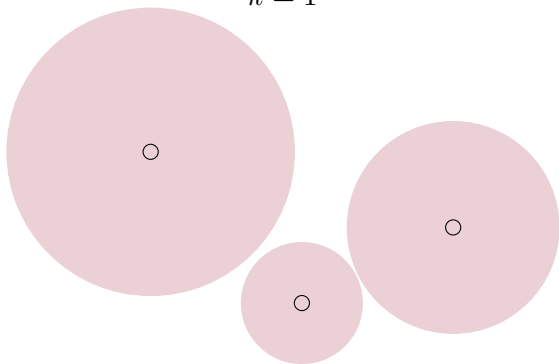
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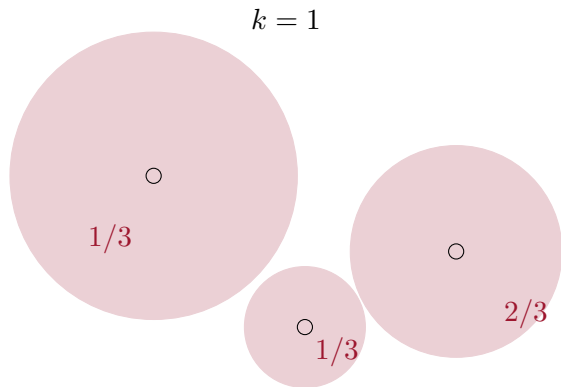
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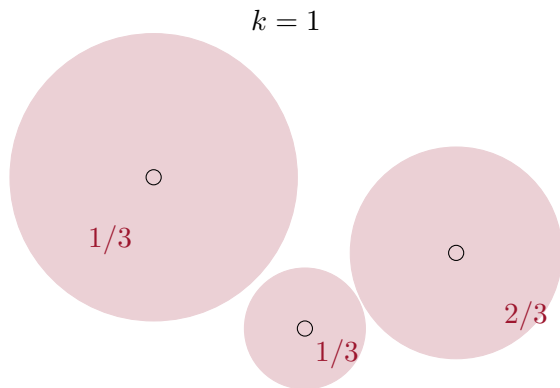




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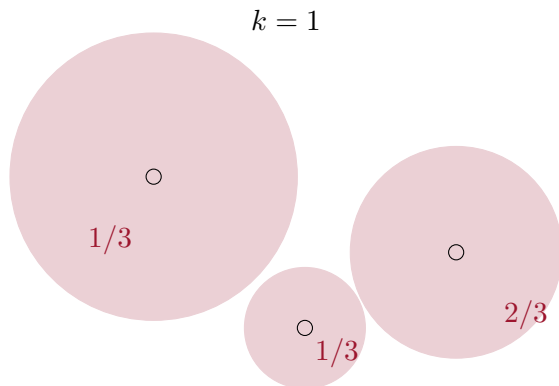


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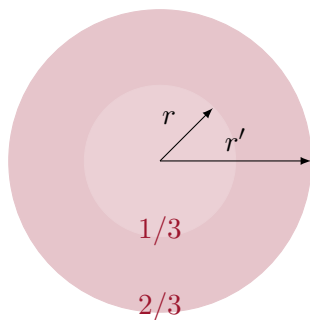


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$$\sum_{B \in \mathcal{B}} y(B) \leq k$$

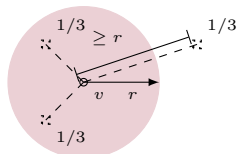
## New LP: Monotonicity constraints

$$y(v, r) \leq y(v, r'), \forall v \in C, r \leq r' \in \mathbb{R}$$

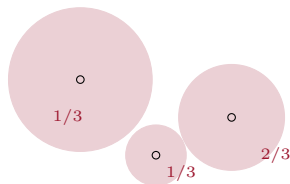


## New LP: Constraints

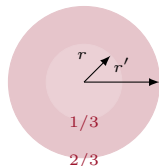
**Consistency.**  $\forall v \in C, r \in \mathbb{R},$   
 $C_v \geq r(1 - y(v, r))$



**Disjoint balls.**  $\forall$  pairwise disj.  $\mathcal{B},$   
 $\sum_{B \in \mathcal{B}} y(B) \leq k$



**Monotonicity.**  $\forall v \in C, r \leq r' \in \mathbb{R},$   
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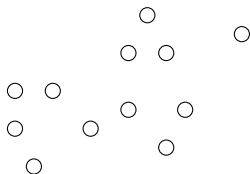
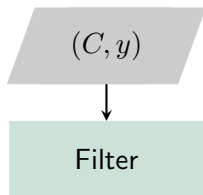
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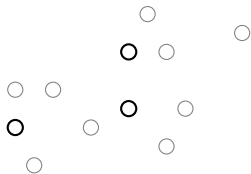
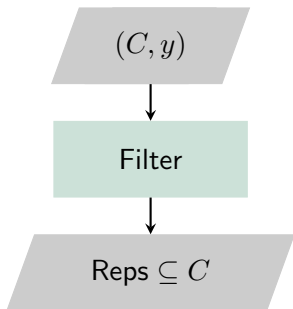
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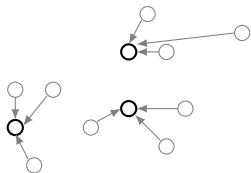
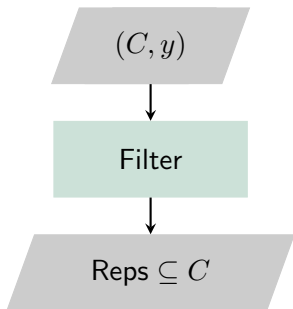




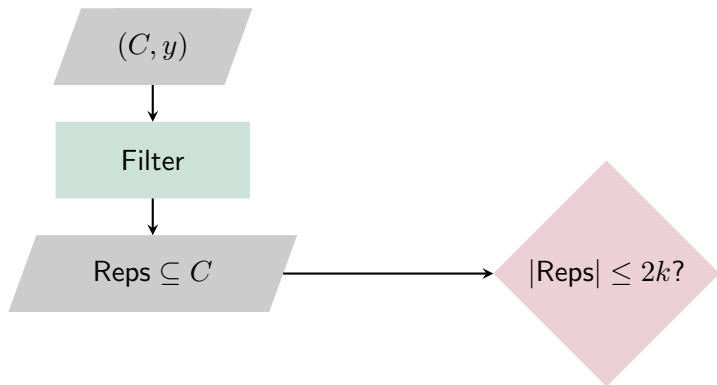
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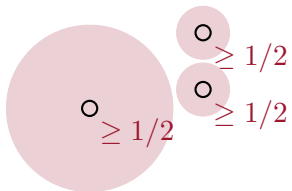
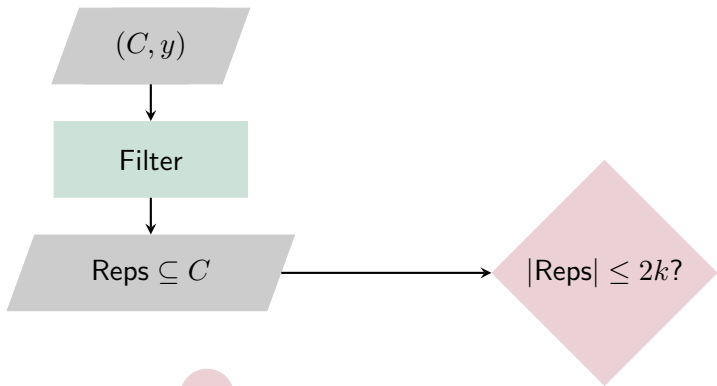
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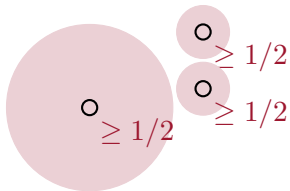
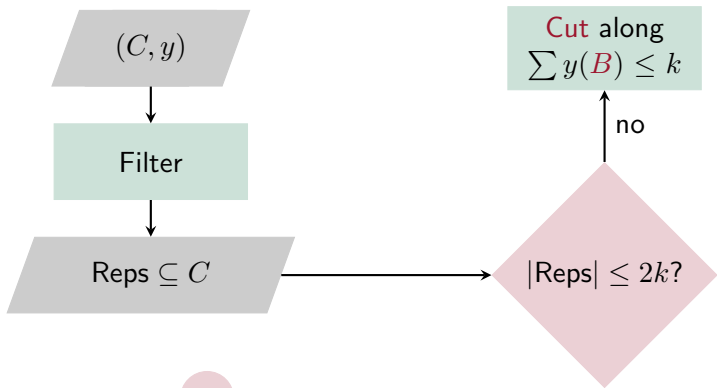
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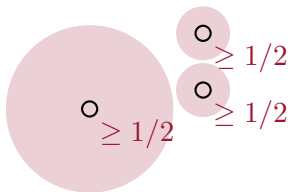
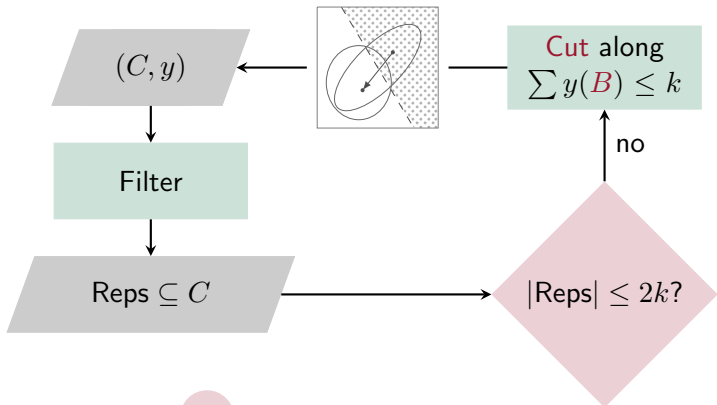
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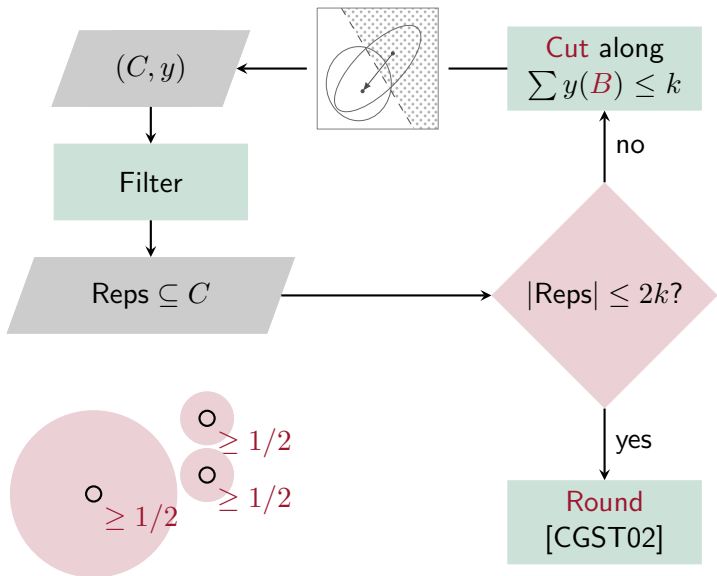
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### Question

Better than 5.35-approx for continuous  $k$ -median?

# Summary

- Continuous vs discrete:  $\beta$  gap?
- New LP:  $C_v, y(v, r)$
- Round-or-cut
- Open: continuous  $k$ -median